

Traffic Analysis Resistant Networks (TARN)

Richard R. Brooks, K.-C. Wang, Lu Yu,
Geddings Barrineau, Jon Oakley and Qing Wang

The Holcombe Department of
Electrical and Computer Engineering
Clemson University,
Clemson, SC, 29634, USA

Email: rrb@g.clemson.edu, kwang@clemson.edu,
lyu@g.clemson.edu, cbarrin@g.clemson.edu,
joakley@clemson.edu, qw@g.clemson.edu

March, 2017

Problem Statement

Network

▷ vulnerabilities

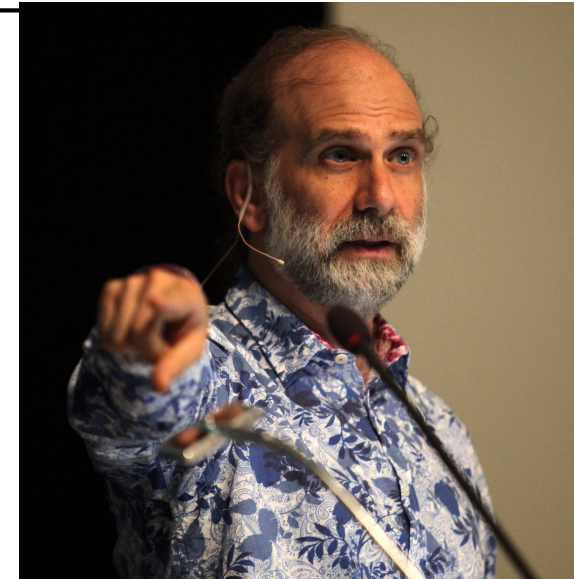
Why does this matter?

Current Solutions

Our Approach

Current Implementation

- Original Internet architecture did not consider security
 - Packets in clear text.
 - Meta-data in clear text.
- SSL/TLS and ssh suite encrypt payloads.
- Meta-data still in clear text.
- IP address meta-data leaks communications partners.
- Port numbers leak protocols.
- Packet sizes identify web content, even when encrypted.
- Timing side-channels leak information (including passwords), even when encrypted.



Why does this matter?



Problem Statement

Network
vulnerabilities

Why does this
▷ matter?

Current Solutions

Our Approach

Current
Implementation

- Censorship
 - DNS filtering
 - IP filtering
- Surveillance
 - Watering hole attacks
 - Identify fans of Breitbart, Democracy Now!, Fox News, MSNBC, WSJ, NYTimes...
- Traffic hijacking
- Man-in-the-middle attacks
- DDoS

Problem Statement

Current Solutions

▷ Existing solutions

What else can be done?

Our Approach

Current Implementation

- VPNs – Closed systems. Typically commercial. Easily detected and blocked. For example, Iran has simply blocked all encrypted connections.
- Single hop proxies – *Psiphon, Lantern, Clemson Internet Freedom in Africa...* Open systems. Subject to blocking of TLS/SSL. Some obfuscate their traffic. Must trust intermediates.
- Onion routing – *Tor* Increased latency and jitter. Access blocked by China and Iran. Bridge nodes get around blocking. Probing finds bridge nodes. Pluggable Transports.
- Decoy routing – *BBN, U Mich, UIUC* Router recognizes encrypted signals in headers and reroutes the traffic. Traffic has to pass through specific routers. Need to trust large multi-national telecom, US Federal Gov., or similar entity.
- Botnets – use DGA and double fast flux to hide their traffic and the location of their servers.

What else can be done?



Problem Statement

Current Solutions

Existing solutions

▷ What else can be done?

Our Approach

Current Implementation

- BGP has issues
 - Fat fingered route leaks
 - Pakistan removed Youtube from global Internet
 - Guadalajara – DC traffic routed through Belarus
 - UK nuclear arms traffic routed through Russia/Ukraine

- Few people monitor BGP route problems and detection takes days.
- IPv6 address space is enormous and almost empty.
- If you map all IPv4 addresses twice onto IPv6, the chance of a collision is essentially zero.
- Could we use Hedy Lamarr's insights on the IP space?

Option 1 – BGP abuse

Problem Statement

Current Solutions

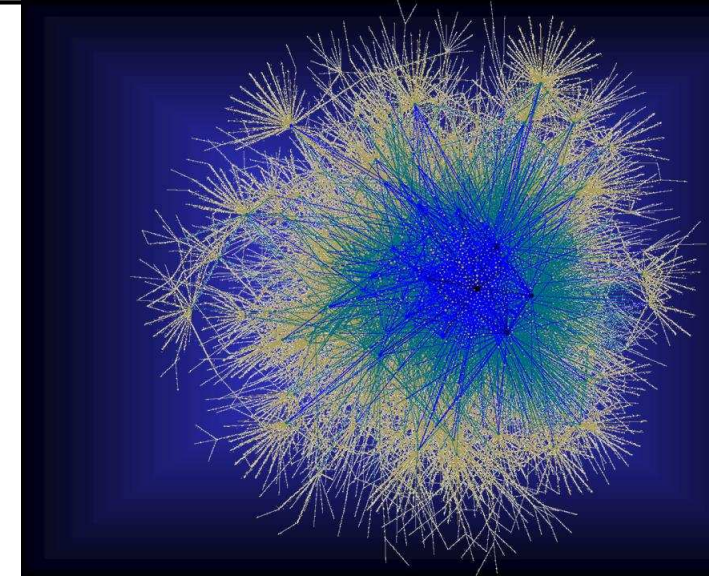
Our Approach

Option 1 – BGP ▷ abuse

Option 2 – SDX

Current Implementation

- We know how to foil DPI, mimic protocols, fool timing side-channels.
- Remaining problem – IP's in clear-text.
- Can we do IP redirection?
- Can we do IP randomization?
- Full randomization limited . BGP squatting violates civil contracts.
- We can test. BGP point of presence changes IP ranges (within allowable ranges) using a schedule shared with the client.
- Intermediate Floodlight Openflow controllers re-write IP destination addresses on the fly.
- Prototype tests concept of IP address hopping.
- Drawbacks: BGP injection overhead, requires synchronization, route update delays, limited number of updates.



Problem Statement

Current Solutions

Our Approach

Option 1 – BGP
abuse

▷ Option 2 – SDX

Current
Implementation

- Next step.
- Instead of BGP, use SDN infrastructure.
- BGP route updates not needed. Software Defined Internet Exchanges (SDX) used.
- Re-routing potentially less transparent.
- Randomization not necessarily end-to-end. Potentially at multiple points.
- Re-routing not necessarily using IP addresses, we have used MAC addresses for similar things.
- Need to establish how to reduce overhead, make more robust, reduce information leakage.
- Lots of possibilities. More questions than answers for now.

Problem Statement

Current Solutions

Our Approach

Current
Implementation

PEERING
▷ Testbed
How We Use
PEERING
Maintaining
Communication
Sessions

- Can an experimenter perform BGP interoperation experiments with today's Internet?
- Yes! The PEERING BGP testbed connects with real networks (via BGP), running experiments to exchange BGP routes and traffic directly with these networks.
- PEERING controls 8 ASNs (e.g. 47065, 61575, etc)
- PEERING controls 3 IPv4 prefixes (e.g. 184.164.224.0/19, 138.185.228.0/22, etc) and 1 IPv6 prefix (e.g. 2804:269c::/32)
- PEERING has points of presence (PoPs) at multiple internet exchange points and universities. Experimenters advertise and withdraw BGP prefixes at selected PoP.



How We Use PEERING

Problem Statement

Current Solutions

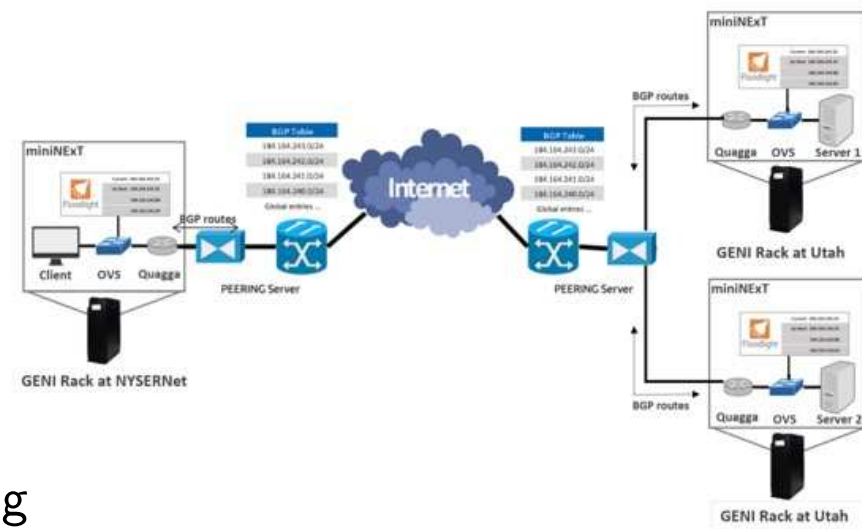
Our Approach

Current Implementation

PEERING Testbed

How We Use
▷ PEERING
Maintaining
Communication
Sessions

- Allotment of four prefixes:
184.164.240-243.0/24
Two PoPs: Amsterdam
and Seattle
- miniNExT network on GENI
node connects to PEERING using
an OpenVPN tunnel
- BGP update limit of 18 update messages per day
- Advertise two type of BGP update messages
 - Announcement
 - Withdrawal
- Swap BGP prefix from one server to another:
 - Withdraw existing BGP prefix on one server
 - Insert that BGP prefix to another server



Maintaining Communication Sessions

Problem Statement

Current Solutions

Our Approach

Current
Implementation

PEERING Testbed

How We Use

PEERING

▷ Maintaining
Communication
Sessions

- Application on host communicates with a TARN server using server's identifier address (e.g. 10.0.0.1)
- Traffic from application travels through an OVS connected to the Floodlight OpenFlow controller
- Floodlight, recognizing the identifier address, inserts flows on the OVS to rewrite the identifier address to the current randomized IP address and vice versa
- New flows are inserted as the randomized IP address changes
- Because the rewrites are transparent to the application, the same socket can be used on both sides of the connection for the duration of the session
- The infrastructure can be located within a single host or on the edge of an

